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# **Real-time Packet Performances under Socket Application**

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*Abstract*— Real-time packet capture and analysis has become a challenging problem, since it plays a vital role not only in identifying the problem causing entity but also to manage the network efficiently. The administrator has to know the network bandwidth and other resources that are used for accounting and auditing. This creates an emphasis to monitor network traffic and conduct analysis to ensure smooth operations performed routinely. Time is an important factor, which contributes to the variations of packet flow. In this paper, initially, the packet movement in the network was monitored with only hypertext transfers using HTTP. The same movement was analyzed with respect to other file transfer applications. The number of file downloads between the HTTP server and clients are also varied to follow the behavior of data. We also designed and developed multithreaded socket applications using java which was again monitored and analyzed using the above methodology. We compared the reading and the experimental results shows that the developed application outperforms in terms of packet acceleration compared to the previous ones.

Keywords- bandwidth, delay, throughput, hypertext, HTTP, server, client, multithread and socket

#### I. INTRODUCTION

Packet analysis could be done by capturing and analyzing traffic passing by the machine where the tool is installed with results displayed. A lot of investigations had been made in bandwidth allocation but gaps in studies presented in the paper show that there is still strong scope for improvement [1]. Whereas Packet analyzer consists of a packet analyzing PC program and sniffer device . The sniffer device of the packet analyzer monitors captures MAC packet frames, and transmits packet frames to the packet analyzing computer [2]. It also decodes all major and frequently used protocols including TCP/IP, UDP, HTTP, etc. In the research work conducted to compare TCP and UDP packet analysis, they concluded that frame length, frame no. and bytes captured during sending a mail more of TCP than UDP [3].

Traffic flow models are based on observations. Traffic flow prediction is an important task for traffic managers and it allows performance assessment of major traffic infrastructure [4]. It is possible to filter the network traffic to focus on the specified needed information. It is flexible and allows powerful filters during or after capture to isolate traffic by specific node, protocol, and packet content. Hence the tool was initialized as per the settings and HTTP server response time was configured as shown in fig 1.

A distinguishing feature of the methods is addressing

problems that need to be solved by interactive systems at a large scale [5]. Network performance and security prevent network problems, conduct effective troubleshooting and take actions quickly to solve possible problems. The network bandwidth and other resources are used for accounting, auditing or for network planning purposes and analysis of the packets passing through the network. Network Packet Analyzer CAPSA is an advanced network traffic monitoring, analysis and reporting tool, based on Windows operating systems. It captures and analyzes all traffic transport over both Ethernet and WLAN networks and decodes all major TCP/IP and application protocols. Its advanced application analysis modules allows to view and log key communication applications such as emails, http traffic, instant messages and DNS queries.

The comprehensive reports and graphic views enable to understand network performance and bandwidth usage quickly, to check network health. Initial HTTP readings were noted by varying the requests as shown in fig 4. Hypertext transfer protocol was designed to be quick, simple, and non instructive. The connection between a server and a client program is temporary and must be reestablished for every data transfer. A URL is always a single, unbroken line of text with no spaces. Web browser generally displays the URL of the Web page currently being viewed near the top of the windows.

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Fig 1: Initial seeting showing buffer size of 512 KB.

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		TCP Connection Initiated				25	
		TCP Connection Established					
		DNS Analysis				Count	
		DNS Request				20	
		DNS Correct Response				12	
		DNS Error Response					
		SMTP Analysis				Count	
		SMTP Connections				0	
		SMTP Messages Sent				0	
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		POP3 Analysis				Count	
		POP3 Connections				0	
		POP3 Messages Received				0	
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		HTTP Connections				24	
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Packet filters:	No filters enabled	Logins				0	
Error packets:	2	Total Messages				0	
Packets captured:	2,147	Messages Sent					
Packets lost:	0	Messages Received				0	
Packets accepted:	2,147	MSN Activities				Count	
Packets rejected:	0	Logins				0	
Buffer usage:	804 KB	Total Messages				N 0	

Fig 2: initial analysis of HTTP at the start

This paper is organaised as Section I contains the introduction of Packet analysis, Section II contain the related work of Packet analysis, Section III contain the methodology which comprises before and after the HTTP requests made without download, Section IV describes results and discussion.

#### **II. RELATED WORK**

Allan M de Souza et al in their paper titled Traffic management systems: A classification, review, challenges, and future perspectives discussed traffic management systems are composed of a set of application and management tools to improve the overall traffic efficiency and safety of the transportation systems. They further added that traffic management system gathers information from heterogeneous sources, exploits such information to identify hazards that may potentially degrade the traffic efficiency, Vol.5(3), June 2017, E-ISSN: 2321-3256

and then provides services to control them. Their article presents a classification, review, challenges, and future perspectives to implement a traffic management system [6].

Vivek Raich et al had discussed that though cloud computing is network of networks Software as a Service (SaaS) model can be seen as a layer with IaaS at the base allowing full control of resources and storage, PaaS in the middle allowing development on an existing platform and finally, SaaS providing limited development opportunities but having appeal to end-users. Though it had limited development oppurtunities the need for software especially in this scenario is essential [7]

### III. METHODOLOGY

The packets used were noted and the readings for corresponding protocols are shown in fig 3 and the initial HTTP requests were shown in fig 4.

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🖶 遭 Broadcast Addresses (1)	Other	2.625 MB	7,387	DFM0 bps	23.887%		
🗉 遭 Multicast Addresses (3)	THTTPS	277.146 KB	540	0 bps	2.463%		
	MSN	DEMO VERSION 121.101 KB	878	0 bps	1.076%		
🛓 🧔 "This" Networks (1)	AIM	14.326 KB	88	0 bps	0.127%		
🚋 🔂 Link Local (2)	T T LDAP	13.465 KB	41	DEMObps	0.120%		
🖃 付 Private-Use Networks (2)	SMTP	7.782 KB	61	0 bps	0.069%		
🗈 付 Multicast Addresses (2)	···· 🚏 Kerberos	DEMO VERSION 7.014 KB	20	0 bps	0.062%		
Broadcast Addresses (1)	🖨 🚏 FTP	4.157 KB	53	0 bps	0.037%		
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	🎁 POP3	3.497 KB	25	0 bps	0.031%		
	🖓 🖓 Yahoo Messen	DEMO VERSION 3.426 KB	32	0 bps	0.030%		
	📮 🚏 NetBIOS	3.348 KB	28	0 bps	0.030%		
	🖓 Session Ser	3.348 KB	28	DEM0 bps	0.030%		
	- 🍯 SCCP	1.965 KB	4	0 bps	0.017%		
	··· 🚏 NFS	DEMO VERSION 1.593 KB	13	0 bps	0.014%		
	🚽 🚏 CIFS	1.136 KB	16	0 bps	0.010%		
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cket filters: No filters enabled	🖨 🍹 UDP	MO VERSION 3.538 MB	17,118	0 bps	32.198%		
ror packets: 0	🖓 Other	2.086 MB	10,121	0 bps	18.988%		
ckets captured: 38,713	P BitTorrent	1.208 MB	5,342	DEM0 bps	10.993%		
ckets lost: 0	📮 🚏 NetBIOS	132.532 KB	880	0 bps	1.178%		
ckets accepted: 38,713	🖓 Datagram S	DEMO VERSION 68.448 KB	287	0 bps	0.608%		
ckets rejected: 0	👘 🎁 Name Servi	64.084 KB	593	0 bps	0.570%		
iffer usage: 13,683 KB		III.			0.50007	Þ	

Fig 3 packets used by protocols

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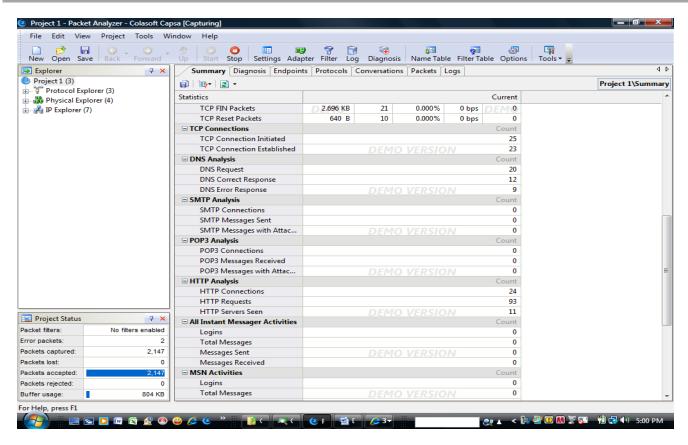


Fig 4 Buffer usage after the HTTP requests made without download

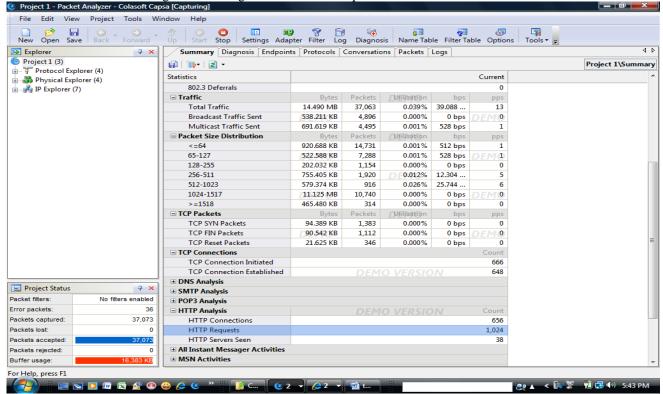


Fig 5 Buffer usage of HTTP with two downloads

The two sites of same URL were requested and the same file of same size was downloaded. Now in order to invoke downloads of same size file single HTTP were noted and two HTTP requests were done for the same URL. The readings were noted for the same file of same size but as different requests as shown in fig 5. From the figure we can view that the buffer usage is almost nearing its capacity indicated in red. This is due to the fact that when more hypertext transfers are made the allocated buffer usage increases.

### IV. RESULTS AND DISCUSSION

#### Results of Socket Application Performance

A socket is a software endpoint that establishes bidirectional communication between a server program and one or more client programs. The socket associates the server program with a specific hardware port on the machine where it runs so any client program anywhere in the network with a socket associated with that same port can communicate with the server program. Typical system configuration places the server on one machine, with the clients on other machines. The clients connect to the server, exchange information, and then disconnect. Moreover, the processes that use a socket can reside on the same system or on different systems on different networks. Initially, client and server create sockets. The server binds to an address and port and it is automatically done in Java. The client knows server's address and port whereas the server listens on that port. The client connects to the address and port server accepts the connection. Hence the client and server read from and write to their sockets.

Client Side Program: import java.io.\*; import java.net.\*; public class myClient { public static void main(String[] args) throws IOException { Socket echoSocket = null; PrintWriter out = null; BufferedReader in = null: try { // connect to local host at port 4444 echoSocket = new Socket ("10.0.0.1", 4444); } The Server side program: import java.net.\*; import java.io.\*; public class myServer { public static void main(String[] throws args) IOException { ServerSocket serverSocket = null; try { serverSocket = new ServerSocket (4444); // binds, listens automatically 

}

}

The application of the client side and server side execution

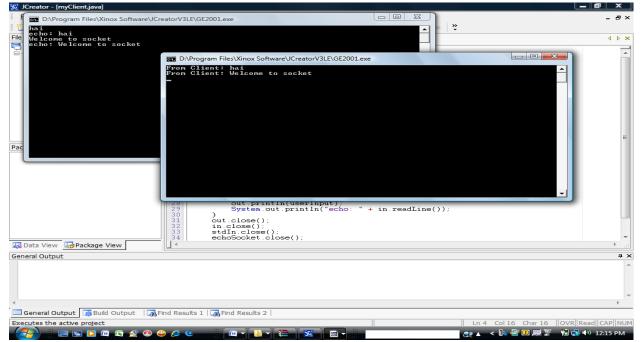


Fig 6 Interaction of server and client programs

Time	Received	Sent	R.Diff	S.Diff	S- R
4:03:00	2,534,091	474,244	178	172	-6
4:04:00	2,534,269	474,429	108	113	5
4:05:00	2,534,377	474,543	136	144	8
4:06:00	2,534,513	474,687	-	-	-
		Total	422	429	

Table 1. Performance measurement of socket application

From the above table the average of sent and received calculated are 2.34 packets per second for receiving and average of sent 2.38 packets per second. Finally, the behavior of the system was compared against the previous reports. Average of received is 2.34 packets per second the average of sent is 2.38 packets per second compared to 2.19 and 2.20 respectively.

#### V. CONCLUSION

The network performance factors delay and throughput were discussed and measured. Real-time packet capture and analysis had been done using packet analysis technique. Also the network comprises Netest which helps in monitoring the bad performing hardware and provides way for better performance. The observation include the load of a large user population among multiple servers were not spread. Reduce network usage and improve server performance when there is a significant user population in distributed locations. Though there were five servers functioning but the configurations are inefficient. The better solution could be the application of same time server which the network lacks currently. The tool had been tested under various HTTP requests and the increased requests are shown. Moreover, the analysis was done with variable number of downloads. The socket application was designed and implemented. The server side and client side communicated effectively and the measurements as well as the performance of this application was also measured repeatedly and shown in table 1. From the above table 1 the average of sent and received calculated are 2.34 packets per second for receiving and average of sent 2.38 packets per second. Finally, the behavior of the system was compared against the previous reports. Average of received is 2.34 packets per second the average of sent is 2.38 packets per second compared to 2.19 and 2.20 respectively. The results obtained shows the application sent and received bytes for the application increases compared to that of the previous throughput.

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